



THE TWO-PEBBLING PROPERTY ON SHADOW GRAPH OF A PATH

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Abstract

Given a distribution of pebbles on the vertices of a connected graph G , a pebbling move is defined as the removal of two pebbles from some vertex and the placement of one of those pebbles at an adjacent vertex. The pebbling number, $f(G)$, of a connected graph G , is the smallest positive integer such that from every placement of $f(G)$ pebbles, we can move a pebble to any specified vertex by a sequence of pebbling moves. A graph G has the 2-pebbling property if for any distribution with more than $2f(G) - q$ pebbles, where q is the number of vertices with at least one pebble, it is possible, using the sequence of pebbling moves, to put 2 pebbles on any vertex. In this paper, we find the pebbling number for the shadow graph of a path and show that it satisfies the 2-pebbling property.

Keywords: Pebbling number, 2-pebbling property, Shadow graph.

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1. Introduction

Pebbling, one of the latest evolutions in graph theory proposed by Lakarias and Saks, has been the topic of vast investigation with significant observations. Having Chung [1] as the forerunner to familiarize pebbling into writings, many other authors too have developed this topic. Hulbert published a survey of graph pebbling [5].

Consider a connected graph with fixed number of pebbles distributed on its vertices. A pebbling move consists of the removal of two pebbles from a vertex and placement of one of those pebbles at an adjacent vertex. The pebbling number of a vertex v in a graph G is the smallest number $f(G, v)$ such that for every placement of $f(G, v)$ pebbles, it is possible to move a pebble to v by a sequence of pebbling moves. Then the pebbling number of G is the smallest number, $f(G)$ such that from any distribution of $f(G)$ pebbles, it is possible to move a pebble to any specified target vertex by a sequence of pebbling moves. Thus $f(G)$ is the maximum value of $f(G, v)$ over all vertices v .

Chung [1] defined the 2-pebbling property of a graph. Given a distribution of pebbles on G , let p be the number of pebbles, q be the number of vertices with at least one pebble, we say that G satisfies the 2 - pebbling property, if it is possible to move two pebbles to any specified vertex whenever p and q satisfy the inequality $p + q > 2f(G)$.

In this paper, we find the pebbling number for the shadow graph of a path and show that the 2-pebbling property. In Section 2, we give some useful results for the subsequent sections. In Section 3, we determine the pebbling number for the shadow

graph of a path $D_2(P_n)$. In Section 4, we prove that the shadow graph of a path $D_2(P_n)$ satisfies the 2-pebbling property.

2. Preliminary

We now introduce some definitions and notations which will be useful for the subsequent sections. For graph theoretic terminologies we refer to [4].

Definition 2.1. The *shadow graph* $D_2(G)$ of a connected graph G is constructed by taking two copies of G , say G_1 and G_2 and joining each vertex u in G_1 to the neighbours of the corresponding vertex v in G_2 .

The shadow graph of a path is denoted by $D_2(P_n)$. Label the vertices in the first copy of the path by x_1, x_2, \dots, x_n and the vertices in the second copy of the path by $x_{n+1}, x_{n+2}, \dots, x_{2n}$ starting from the left.

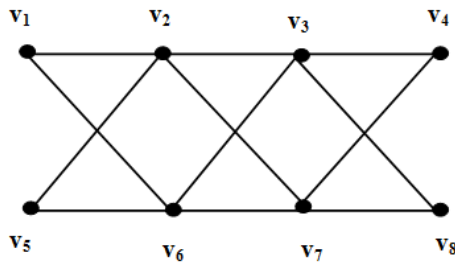


Figure 1.1. $D_2(P_n)$

Theorem 2.2 [3] Let P_n be a path on n vertices. Then $f(P_n) = 2^{n-1}$.

Theorem 2.3. [2] Let $K_{1,n}$ be a star graph, where $n > 1$. Then $f(K_{1,n}) = n + 2$.

Theorem 2.4. [6] (i). Let C_{2k} be an even cycle on $2k$ vertices. Then $f(C_{2k}) = 2^k$.

(ii) Let C_{2k+1} be an odd cycle on $2k + 1$ vertices. Then

$$f(C_{2k+1}) = 2 \left\lfloor \frac{2^{k+1}}{3} \right\rfloor + 1.$$

Theorem 2.5. [6] Let G be a graph with diameter $G = 2$. Then G has the 2-pebbling property.

Theorem 2.6. [3] All paths satisfy the 2-pebbling property.

Theorem 2.7. [3] All cycles have the 2-pebbling property.

3. Pebbling on the Shadow graph of a path $D_2(P_n)$

Remark 3.1. A distribution of pebbles on the vertices of the graph G is a function $p : V(G) \rightarrow N \cup \{0\}$. Let $p(v)$ denote the number of pebbles on the vertex v and $p(A)$ denote the number of pebbles on the vertices of the set $A \subseteq V(G)$. Let v be a target vertex in the graph G . If $p(v) = 1$ or $p(u) \geq 2$, where $uv \in E(G)$, then we can move a pebble to v easily. So we always assume that $p(v) = 0$ and $p(u) \leq 1$ for all $uv \in E(G)$, when v is the target vertex.

For $n = 2$, $D_2(P_2)$ is isomorphic to C_4 , we have the following theorem.

Theorem 3.2. [3] For the shadow graph of the path P_2 , $f(D_2(P_2)) = 4$.

Theorem 3.3. For the shadow graph of the path P_3 , $f(D_2(P_3)) = 6$.

Proof. Placing 3 pebbles on x_1 and one pebble on each x_3 and x_4 , we cannot reach x_6 . Thus $f(D_2(P_3)) \geq 6$.

Now we prove that $f(D_2(P_3)) \leq 6$. Let D be any distribution of 6 pebbles on the vertices of $D_2(P_3)$.

Case 1 : Let x_3 be the target vertex.

Clearly $p(x_3) = 0$, $p(x_2) \leq 1$ and $p(x_6) \leq 1$ by Remark 3.1. If $p(x_2) = 1$ we can move another pebble to x_2 , since x_1 or x_4 or x_6 contains at least 2 pebbles. Therefore assume that $p(x_2) = 0$. Since $\langle D_2(P_3) - \{x_2\} \rangle$ is isomorphic to $K_{1,4}$ and $f(K_{1,4}) = 6$, we are done. Similarly, we are done if x_1 , x_4 and x_6 are the target vertices.

Case 2 : Let x_2 be the target vertex.

Clearly $p(x_2) = 0$, $p(x_1) \leq 1$, $p(x_3) \leq 1$, $p(x_4) \leq 1$ and $p(x_6) \leq 1$. Suppose $p(x_i) = 1$ for some $i \in \{1, 3, 4, 6\}$. Then using the path $P : x_5 x_i x_2$ we can pebble the target. Suppose $p(x_i) = 0$ for every $i \in \{1, 3, 4, 6\}$. Then $p(x_5) \geq 5$ and hence we are done. Similarly, we are done if x_5 is the target vertex.

Theorem 3.4. *For the shadow graph of a path P_n ,*

$$f(D_2(P_n)) = 2^{n-1} + 2, (n \geq 4).$$

Proof. Placing 2^{n-1} pebbles on x_1 and one pebble on each x_n and x_{n+1} , we cannot reach x_{2n} . Thus $f(D_2(P_n)) \geq 2^{n-1} + 2$. Now we prove the sufficient part by induction on n . Clearly, it is true for $n = 3$, by Theorem 3.3. So, we assume the result is true for $4 \leq n' < n$. Let D be any distribution of $2^{n-1} + 2$ pebbles on the vertices of $D_2(P_n)$.

Case 1 : Let v be any target vertex other than x_1, x_n, x_{n+1} and x_{2n} .

Clearly, $p(v) = 0$ by Remark 3.1. We note that $d(u, x_1) \leq n - 2$ and $d(u, x_{n+1}) \leq n - 2$ for every $u \notin \{x_n, x_{2n}\}$. If $p(x_1) \geq 2^{n-2}$ or $p(x_{n+1}) \geq 2^{n-2}$, then we are done. Therefore assume that $p(x_1) \leq 2^{n-2} - 1$ and $p(x_{n+1}) \leq 2^{n-2} - 1$. By moving as many pebbles as possible from x_1 to x_2 and from x_{n+1} to x_{n+2} , we see that the subgraph $\langle D_2(P_n) - \{x_1, x_{n+1}\} \rangle \cong D_2(P_{n-1})$ contains at least $2^{n-2} + 2$ pebbles and hence we are done, by induction.

Case 2 : Let v be any target vertex, where $v \in \{x_1, x_n, x_{n+1}, x_{2n}\}$.

Without loss of generality, we assume that $v = x_1$. Let $p(x_2) + p(x_3) + \dots + p(x_{n-1}) + p(x_{n+1}) + \dots + p(x_{2n-1}) = p$. Suppose $p \geq 3$. Then $p(x_n) + p(x_{2n}) = 2^{n-1} + 2 - p \leq 2^{n-1} - 1$. After

moving as many pebbles as possible from x_n to x_{n-1} and from x_{2n} to x_{2n-1} , we see that $\langle D_2(P_n) - \{x_n, x_{2n}\} \rangle \cong D_2(P_{n-1})$ contains at least $2^{n-2} + 2$ pebbles and hence we are done, by induction. Therefore assume that $p \leq 2$.

Subcase 2.1: Let $p = 0$.

Clearly, $p(x_n) + p(x_{2n}) = 2^{n-1} + 2$. After moving as many pebbles as possible from x_n to x_{n-1} and from x_{2n} to x_{n-1} , the vertex x_{n-1} contains at least 2^{n-2} pebbles. Since $d(x_1, x_{n-1}) = n - 2$, we are done.

Subcase 2.2: Let $p = 1$.

Now $p(x_n) + p(x_{2n}) = 2^{n-1} + 1$. After moving as many pebbles as possible from x_n to x_{n-1} and from x_{2n} to x_{n-1} , the vertex x_{n-1} contains at least $\left\lfloor \frac{2^{n-1} + 1 - 1}{2} \right\rfloor = 2^{n-2}$ pebbles. Since $d(x_1, x_{n-1}) = n - 2$, we are done.

Subcase 2.3: Let $p = 2$.

Now $p(x_n) + p(x_{2n}) = 2^{n-1}$. Then we can move at least $\left\lfloor \frac{2^{n-1} - 2}{2} \right\rfloor = 2^{n-2} - 1$ pebbles to x_{n-1} or x_{2n-1} . Hence we can reach the vertex x_2 . If $p(x_{n+1}) = 2$ we can move one pebble to x_2 and hence we reach the target. Therefore assume that $p(x_{n+1}) \leq 1$. Since $p = 2$, there exists an j such that the vertex x_j is occupied, where $j \in \{2, 3, \dots, n-1, n+2, \dots, 2n-1\}$ and hence we can easily reach the target.

4. The 2-pebbling property

In this section, we show that the shadow graph of a path $D_2(P_n)$ satisfies the 2-pebbling property. Since $D_2(P_2)$ is isomorphic to C_4 and by Theorem 2.7, $D_2(P_2)$ has the 2-pebbling property.

Remark 4.1. Consider the graph G with n vertices and $2f(G) - q + 1$ pebbles on it and we choose a target vertex v from G . If $p(v) = 1$, then the number of pebbles remaining in G is $2f(G) - q \geq f(G)$, since $f(G) \geq n$ and $q \leq n$, and hence we can move the second pebble to v . Let us assume that $p(v) = 0$. If $p(u) \geq 2$, where $uv \in E(G)$, we move a pebble to v from u . Then the graph G has at least $2f(G) - q + 1 - 2$ pebbles, since $f(G) \geq n$ and $q \leq n - 1$, and hence we can move the second pebble to v . So, we always assume that $p(v) = 0$ and $p(u) \leq 1$ for all $uv \in E(G)$, when v is the target vertex.

Theorem 4.2. *The graph $D_2(P_3)$ satisfies the two-pebbling property.*

Proof. Since the diameter of $D_2(P_3)$ is two, by Theorem 2.5 we conclude that, the graph $D_2(P_3)$ satisfies the two-pebbling property.

We first prove the following lemma.

Lemma 4.3. *Given any distribution of $2f(D_2(P_n)) - q + 1$ pebbles on the vertices of $D_2(P_n)$, $n \geq 3$, we can move two pebbles to x_1 and retain at least two*

pebbles on $D_2(P_n)$. If $p(x_{n+1}) = 0$ we can move two pebbles to x_1 and retain at least 3 pebbles on $D_2(P_n)$.

Proof. The proof is by induction on n . Let $n = 3$. Let D be any distribution of $2(6) - q + 1 = 13 - q$ pebbles on it. Clearly $p(x_1) = 0$, $p(x_2) \leq 1$ and $p(x_5) \leq 1$ by Remark 4.1. Then $q \leq 5$.

Suppose $q = 5$. Now 8 pebbles are distributed on the vertices of the graph $D_2(P_3)$. Also $p(x_1) = 0$, $p(x_2) = 1$ and $p(x_5) = 1$. Thus we can easily move 2 pebbles to x_1 using 6 pebbles and hence we can retain 2 pebbles. Therefore assume that $q \leq 4$. Then either x_2 or x_5 is occupied.

Without loss of generality, assume $p(x_2) = 1$. Then using 2 pebbles we can move an additional pebble to x_2 and hence we can move a pebble to x_1 . Now the remaining number of pebbles distributed on the vertices of $D_2(P_3)$ is at least $13 - q - 3 \geq 10 - q \geq 6$, since $q \leq 4$. By Theorem 2.3, we can move an additional pebble to x_1 at a cost of at most $2^2 = 4$ pebbles and hence we can retain 2 pebbles. Similarly we are done, if $p(x_2) = 0$ and $p(x_5) = 1$. Therefore assume that $p(x_2) = 0 = p(x_5)$.

If $p(x_4) \geq 2$, then one pebble can be moved to x_2 . Using 2 pebbles we can move an additional pebble to x_2 and thus we can move a pebble to x_1 . Now the remaining number of pebbles distributed on the vertices of $D_2(P_3)$ is at least $13 - q - 4 = 9 - q \geq 6$, since

$q \leq 3$. Then by Theorem 3.3, we can move a pebble to x_1 at a cost of at most 4 pebbles and hence we can retain at least 2 pebbles. Therefore assume that $p(x_4) \leq 1$.

Suppose $p(x_4) = 0$. Then $p(x_3) + p(x_6) = 13 - q \geq 11$. Clearly using 8 pebbles we can move 2 pebbles to x_1 and retain at least 3 pebbles. Therefore assume $p(x_4) = 1$. Then $p(x_3) + p(x_6) = 13 - q \geq 10$. Again using 8 pebbles we can move 2 pebbles to x_1 and retain at least 2 pebbles. Hence the lemma holds, when $n = 3$.

Assume the Lemma is true for $4 \leq n' < n$. Let D be any distribution of $2(2^{n-1} + 2) - q + 1$ pebbles on the vertices of $D_2(P_n)$. Clearly, $p(x_1) = 0$, $p(x_2) \leq 1$ and $p(x_{n+2}) \leq 1$ by Remark 4.1. Suppose $q = 2n - 1$. Then $p(x_2) = 1$ and $p(x_{n+2}) = 1$. Using $2n$ pebbles, we can easily put 2 pebbles to x_1 and the remaining number of pebbles is

$$2(2^{n-1} + 2) - q + 1 - 2n = 2^n + 6 - 4n = 2 + (2^n + 4 - 4n).$$

Since $n \geq 4$, $2(2^{n-1} + 2) - q + 1 - 2n \geq 2$. Hence we can retain at least 2 pebbles on $D_2(P_n)$. Therefore assume that $q \leq 2n - 2$. Then either x_2 or x_{n+2} is occupied. Without loss of generality, assume $p(x_2) = 1$ and $p(x_{n+2}) = 0$. We can move a pebble to x_2 at a cost of at most 2^{n-2} pebbles and hence we can move a pebble to x_1 . Now the remaining number of pebbles is at least

$$2(2^{n-1} + 2) - q + 1 - 2^{n-2} - 1 = 2^{n-1} + 2 + 2^{n-2} + 2 - q.$$

Since $q \leq 2n - 2$ and $n \geq 4$, we conclude that

$$2(2^{n-1} + 2) - q + 1 - 2^{n-2} - 1 \geq 2^{n-1} + 2 = f(D_2(P_n)).$$

By Theorem 3.4, we can move a pebble to x_1 at a cost of at most 2^{n-1} pebbles and retain at least 2 pebbles. Similarly we are done, if $p(x_2) = 0$ and $p(x_{n+2}) = 1$. Therefore assume $p(x_2) = 0$ and $p(x_{n+2}) = 0$. If $p(x_{n+1}) \geq 2$, one pebble can be moved to x_2 . And we can move an additional pebble to x_2 at a cost of at most 2^{n-2} pebbles, since $dist(x_2, u) \leq n - 2$, for every $u \in V(D_2(P_n))$. Thus we can put 2 pebbles to x_2 and hence a pebble to x_1 . Now the remaining number of pebbles is at least

$$2(2^{n-1} + 2) - q + 1 - 2^{n-2} - 2 \geq 2^{n-1} + 2 = f(D_2(P_n)),$$

since $q \leq 2n - 3$. By Theorem 3.4, we can move a pebble to x_1 at a cost of at most 2^{n-1} pebbles and retain at least 2 pebbles. Therefore assume $p(x_{n+1}) \leq 1$.

Case 1: Let $p(x_{n+1}) = 0$.

In this case, $2(2^{n-1} + 2) - q + 1$ pebbles are distributed on the vertices of $\langle D_2(P_n) - \{x_1, x_2, x_{n+1}, x_{n+2}\} \rangle$. Thus using $2(2^{n-2} + 2) - q + 1$ pebbles, we can move 2 pebbles to x_2 and retain at least 3 pebbles,

since $p(x_{n+2}) = 0$ and by induction. Therefore one pebble can be moved to x_1 . Now the remaining number of pebbles is at least

$$2(2^{n-1} + 2) - q + 1 - [2(2^{n-2} + 2) - q + 1] + 3 = 2^{n-1} + 2 + 1 = f(D_2(P_n)) + 1.$$

By Theorem 3.4, we can move a pebble to x_1 at a cost of at most 2^{n-1} pebbles and retain at least 3 pebbles.

Case 2: Let $p(x_{n+1}) = 1$.

Now using $2(2^{n-2} + 2) - (q - 1) + 1$ pebbles, we can move 2 pebbles to x_2 and retain at least 3 pebbles, since $p(x_{n+2}) = 0$ and by induction. Therefore one pebble can be moved to x_1 . Now the remaining number of pebbles is at least

$$2(2^{n-1} + 2) - q + 1 - [2(2^{n-2} + 2) - (q - 1) + 1] + 3 = 2^{n-1} + 2 = f(D_2(P_n)).$$

By Theorem 3.4, we can move a pebble to x_1 at a cost of at most 2^{n-1} pebbles and retain at least 2 pebbles on $D_2(P_n)$.

Theorem 4.4. *The graph $D_2(P_n)$ satisfies the two-pebbling property.*

Proof. The proof is by induction on n . Clearly, the result is true for $n = 3$ by Theorem 4.2. Assume the theorem is true for $4 \leq n' < n$.

Let D be any distribution of $2f(D_2(P_n)) - q + 1 = 2(2^{n-1} + 2) - q + 1$ pebbles on the vertices of $D_2(P_n)$.

Case 1: Let x_1 be the target vertex.

By Lemma 4.3, we can move 2 pebbles to x_1 . By symmetry we are done, if x_n, x_{n+1} and x_{2n} are the target vertices.

Case 2 : Let x_2 be the target vertex.

Clearly $p(x_2) = 0$, $p(x_1) \leq 1$, $p(x_3) \leq 1$, $p(x_{n+1}) \leq 1$ and $p(x_{n+2}) \leq 1$ by Remark 4.1. Now the number of pebbles distributed on the vertices of $\langle D_2(P_n) - \{x_1, x_{n+1}\} \rangle$ is at least $2(2^{n-1} + 2) - q + 1 - 2 \geq 2(2^{n-2} + 2) - q_1 + 1$, where q_1 is the number of occupied vertices in $\langle D_2(P_n) - \{x_1, x_{n+1}\} \rangle$. Clearly $q \geq q_1$. Thus we can move 2 pebbles to x_2 by induction. By symmetry we are done, if x_{n-1}, x_{n+2} and x_{2n-1} are the target vertices.

Case 3 : Let x be the target vertex, other than $x_1, x_2,$

$x_{n-1}, x_n, x_{n+1}, x_{n+2}, x_{2n-1}$ and x_{2n} .

If $p(x_1) + p(x_{n+1}) \geq 2^{n-2} + 1$, then we can move as many pebbles as possible from x_1 and x_{n+1} to x_2 and hence by using the path $P: x_2 x_3 \cdots x$, we are done.

Suppose $p(x_1) + p(x_{n+1}) \leq 2^{n-2}$. Then $\langle D_2(P_n) - \{x_1, x_{n+1}\} \rangle$ has at least $2(2^{n-1} + 2) - q + 1 - 2^{n-2} \geq 2(2^{n-2} + 2) - q_1 + 1$, where q_1 is the number of occupied vertices in $\langle D_2(P_n) - \{x_1, x_{n+1}\} \rangle$, since $q > q_1$. Thus we can move 2 pebbles to x by induction.

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